

CLAIMS

Having thus described our invention, what we claim as new, and desire to secure by Letters Patent is:

1 Sub 1. A method for mapping a valid stack up to a destination program counter,
2 A6 comprising:
3 mapping a path of control flow on the stack from any start point in a
4 selected method to the destination program counter; and
5 simulating stack actions for executing bytecodes along said path.

□ 1 2. The method of claim 1 wherein the step of mapping a path of control flow on the
stack comprises:
3 processing a first linear bytecode sequence until the control flow is interrupted;
4 and recording unprocessed targets from any branches in the first linear bytecode
5 sequence for future processing.

Sub 1 2. The method of claim 2 wherein the step of mapping a path of control flow on the
stack further comprises:
3 processing an additional bytecode linear sequence until the control flow is
4 interrupted; and
5 recording unprocessed targets from any branches in the additional linear bytecode
6 sequence for future processing, where the destination program counter was not
7 reached during an earlier processing of a linear bytecode sequence.

1 2. The method of claim 2 wherein the step of processing any linear bytecode
sequence comprises:

3 determining if a bytecode in said any linear bytecode sequence is a breakpoint
4 with a pointer to bytecode data; and
5 replacing the breakpoint with the bytecode data.

1 5. The method of claim 3 wherein the step of processing any linear bytecode
2 sequence comprises:

3 determining if a bytecode in said any linear bytecode sequence is a breakpoint
4 with a pointer to bytecode data; and
5 replacing the breakpoint with the bytecode data.

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6. The method of claim 1 wherein the step of simulating stack actions executing the
bytecodes along the path further comprises generating a virtual stack.

1 7. The method of claim 6, further comprising:

2 encoding the virtual stack as a bitstring and storing the bitstring at a selected
3 destination for use in memory management operations.

1 8. The method of claim 7, wherein the step of storing the bitstring comprises storing
2 the bitstring to the selected method as compiled on a heap.

1 9. The method of claim 7, wherein the step of storing the bitstring comprises storing
2 the bitstring to a pre-allocated area on the stack.

1 10. The method of claim 1 wherein the step of simulating stack actions executing the
2 bytecodes along the path further comprises:

3 inserting pre-determined stack actions for bytecodes maintaining the control flow
4 in the selected method; and

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calculating stack actions for bytecodes transferring the control flow from the selected method.

1 Sub A7 11. A method for mapping a Java bytecode stack up to a destination program counter comprising:

3 mapping a path of control flow on the stack from any start point in a selected
4 method to the destination program counter; and
5 simulating stack actions for executing bytecodes along said path.

1 Sub A7 12. The method of claim 11 wherein the step of mapping a path of control flow on the stack comprises:
3 processing a first linear bytecode sequence until the control flow is interrupted;
4 and recording unprocessed targets from any branches in the first linear bytecode sequence for future processing.

1 Sub A7 13. The method of claim 12 wherein the step of mapping a path of control flow on the stack further comprises:
3 processing an additional bytecode linear sequence until the control flow is interrupted; and
5 recording unprocessed targets from any branches in the additional linear bytecode sequence for future processing, where the destination program counter was not reached during an earlier processing of a linear bytecode sequence.

1 Sub A7 14. The method of claim 12 wherein the step of processing any linear bytecode sequence comprises:
3 determining if a bytecode in said any linear bytecode sequence is a breakpoint with a pointer to bytecode data; and

5 replacing the breakpoint with the bytecode data.

1 15. The method of claim 13 wherein the step of processing any linear bytecode
2 sequence comprises:

3 determining if a bytecode in said any linear bytecode sequence is a breakpoint
4 with a pointer to bytecode data; and

5 replacing the breakpoint with the bytecode data.

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16. The method of claim 11 wherein the step of simulating stack actions executing the
2 bytecodes along the path further comprises generating a virtual stack.

□ 1 17. The method of claim 16 further comprising:

□ 2 encoding the virtual stack as a bitstring and storing the bitstring at a selected
3 destination for use in memory management operations.

□ 1 18. The method of claim 17, wherein the step of storing the bitstring comprises
2 storing the bitstring to the selected method as compiled on a heap.

□ 1 19. The method of claim 17, wherein the step of storing the bitstring comprises
2 storing the bitstring to a pre-allocated area on the stack.

□ 1 20. The method of claim 11 wherein the step of simulating stack actions executing the
2 bytecodes along the path further comprises:

□ 3 inserting pre-determined stack actions for bytecodes maintaining the control flow
4 in the selected method; and

□ 5 calculating stack actions for bytecodes transferring the control flow from the
6 selected method.

1 21. A computer-readable memory for storing the instructions for use in the execution
2 in a computer of the method of claim 1.

1 22. A computer readable memory for storing the instructions for use in the execution
2 in a computer of the method of claim 11.

3 23. A program storage device readable by a machine, tangibly embodying a program
4 of instructions executable by the machine to perform method steps for mapping a
5 valid stack up to a destination program counter, said method steps comprising:

6 mapping a path of control flow on the stack from any start point in a selected
7 method to the destination program counter; and

8 simulating stack actions for executing bytecodes along said path,

9 wherein the step of mapping a path of control flow on the stack comprises:

10 processing a first linear bytecode sequence until the control flow is interrupted;

11 and

12 recording unprocessed targets from any branches in the first linear bytecode
13 sequence for future processing, and

14 where the destination program counter was not reached during an earlier
15 processing of a linear bytecode sequence,

16 processing an additional bytecode linear sequence until the control flow is
17 interrupted; and

18 recording unprocessed targets from any branches in the additional linear bytecode
19 sequence for future processing.